

MA103 - Class 5

1 Let R be an equivalence relation on a set S . Prove that the following properties hold.

(a) For all $x, y \in S$ we have $xRy \iff [x] = [y]$.

(b) For all $x, y \in S$ we have $\neg xRy \iff [x] \cap [y] = \emptyset$.

a) The statement is: $\forall x, y \in S, (xRy \implies [x] = [y])$
not $(\forall x, y \in S, xRy) \implies [x] = [y]$, which has no meaning.

" \implies " Let $x, y \in S$ arbitrary such that xRy . Since R is transitive, $\forall a \in S, aRx \implies aRy$ (since by hypothesis xRy).
Therefore, let $z \in [x]$. By definition of $[x]$, we have zRx , and therefore zRy . This implies that $\forall a \in [x], a \in [y]$.

" \impliedby " Let $x, y \in S$ such that $[x] = [y]$. Since R is reflexive, we have xRx and therefore $x \in [x] = [y]$. By definition, this means xRy .

b) How to prove an implication by contradiction: Say we want to prove $P \implies Q$ by contradiction. Then we assume that P AND $\neg Q$ hold, and we find a contradiction.

If you want to prove by contradiction that $\forall x \in A, P(x) \implies Q(x)$ you assume $\exists x \in A, P(x) \wedge \neg Q(x)$.

" \implies " Assume by contradiction that there exist $x, y \in S$ with $\neg xRy$ and $[x] \cap [y] \neq \emptyset$. Let $z \in [x] \cap [y]$ (which we can do since $[x] \cap [y] \neq \emptyset$). By definition, xRz and yRz . Since R is symmetric and transitive, this gives us xRy .

" \impliedby " Assume by contradiction there exist x, y such that $\neg (\neg xRy)$ and $[x] \cap [y] = \emptyset$. We know that $[x] = [y]$ by the previous exercise. Moreover, $[x] \neq \emptyset$ since R is reflexive and therefore $x \in [x]$.

2 In lectures, we gave a construction for the rational numbers. This started by looking at the set S of all pairs of the form (a, b) , with $a, b \in \mathbb{Z}$ and $b \neq 0$, and then considering the relation Q on $S = \mathbb{Z} \times (\mathbb{Z} \setminus \{0\})$ defined by:

$$(a, b)Q(c, d) \quad \text{if and only if} \quad ad = bc.$$

(a) Explain why things would go badly wrong if we allow S to include pairs (a, b) with $b = 0$.

(Hint: the answer does not involve 'division by zero'.)

We defined the set Q to be the set of equivalence classes of the relation Q , and defined an "addition" operation on Q by setting $[(a, b)] \oplus [(e, f)] = [(af + be, bf)]$, for each $(a, b), (e, f) \in S$.

• You CANNOT get a contradiction from a definition.

The idea is to show that one of the needed properties does not work.

$$(3, 2) R (0, 0); \quad (0, 0) R (1, 1)$$

but $\neg (3, 2) R (1, 1).$ (calculations needed).

Now we "define" an operation. We just saw that we cannot just "define" stuff. We need to check that everything works properly.

$$e.g. \quad f(x) = \begin{cases} x^2 & \text{if } x^2 - 2 \geq 0 \\ x^2 + 1 & \text{if } x^3 - 5x^2 + \sqrt{3}x \geq 0 \end{cases}$$

what could go wrong in this example? What do we need to check?
 What is the equivalent result in our problem?

We need to check that if $(a, b)Q(a', b')$ and $(c, d)Q(c', d')$ then $[(a, b)] \oplus [(c, d)] = [(a', b')] \oplus [(c', d')]$.

(b) Suppose that $(a, b), (c, d)$ and (e, f) are in S , and that $(a, b)Q(c, d)$. Show that $(af + be, bf)Q(cf + de, df)$.

Your answer should *only* talk about operations with integers. If your answer involves writing any fractions at any stage, it is wrong.

Use this to show that if $(r, s)Q(t, u)$ and $(v, w)Q(x, y)$, then

$$[(r, s)] \oplus [(v, w)] = [(t, u)] \oplus [(x, y)].$$

(This means that the addition operation defined on Q is *well-defined*.)

we know $ad = bc$, and we want to prove $(af + bc)df = (cf + dc)bf$

CLASS SLIDES for WEEK 6.